# Verification of Featherweight Java Programs via Transformation to Higher-order Functional Programs with Recursive Data Types

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### 1. Our ultimate goal

Precise and fully-automated verification of Java programs

This poster :

Proposes a precise and semi-automated verification method for assertion safety of programs in **Featherweight Java (FJ) [Igarashi et al. '01]** extended with **booleans, integers,** 

### 3. Translation from FJ to ML

(This translation is inspired by [Kobayashi and Igarashi '13] ) FJ program  $P_{list} \begin{array}{c|c} Simulate dynamic dispatch using higher$  $order functions and recursive data types \\\hline (Conjecture) \\FJ program \not\rightarrow^*_{FJ} assert false \Leftrightarrow \\ML program \not\rightarrow^*_{ML} assert false \\\hline \end{array}$ 

#### conditional branches, and assertions

## 2. Challenges

Assertion safety verification for ext. FJ often requires

- Context-sensitive analysis of dynamic dispatch
- Path-sensitive analysis of conditional branches

```
class List {
    Int length() { return assert false; }

FJ program P<sub>list</sub>
```

```
class Nil extends List {
    Int length() { return 0; }
```

```
class Cons extends List {
```

```
Int hd; List tl;
Cons(Int hd, List tl) { this.hd=hd; this.tl=tl; }
```

Int length() {
 return 1 + this.tl.length();

Behaves differently depending on the context of length

type obj = Obj of (int -> obj-> obj) \* (obj -> bool)

Encode objects as recursive data structures

let length\_Main this = assert false

let mk\_n\_list\_Main n this = ...

let send\_mk\_n\_list (Obj (m1,m2)) = m1 (Obj(m1,m2))

To simulate dynamic dispatch, extract and call a function

••• from the recursive data structure that encodes an object

```
let main () =
   assert (send_length
      send mk n
```

```
send_mk_n_list_Main
  (Obj(mk_n_list_Main 3, length_Main)) = 3)
```

## 4. Limitation of RCaml and our solution

#### Limitation

RCaml **cannot** handle recursive data structures **context-sensitively** if they contain **functions** (otherwise, it can)

#### Solution

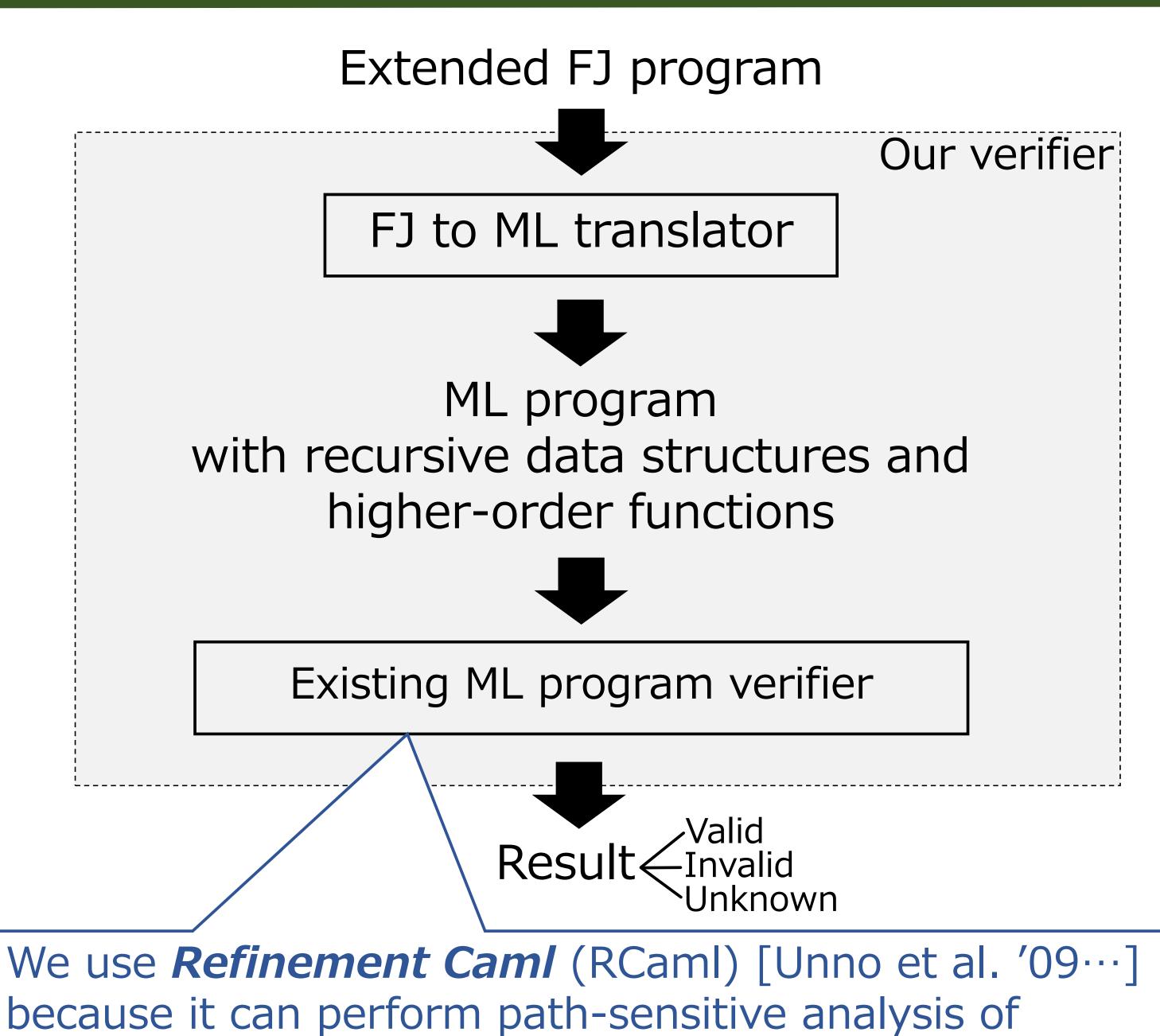
Extend the FJ-to-ML translation to insert context information to the recursive data structures that encode objects

type obj = Obj of cls \* (int -> obj-> obj) \* cls \* (obj -> bool)

class Main {
 List mk\_n\_list(Int n) {
 return if n=0 then new Nil()
 else new Cons(n, this.mk\_n\_list(n-1));
 }
 Behaves differently depending on the run time value of n
}
main() {
 return assert (new Main().mk\_n\_list(3).length() = 3);

## 2. Approach

conditional branches



```
and cls = Mk_n_list_Main of int | Length_Cons | ...
let send_mk_n_list (Obj (cm1,m1,cm2,m2)) = m2 (Obj(cm1,m1,cm2,m2))
...
let main () =
  assert (send_length
      send_mk_n_list_Main
      (Obj(Mk_n_list_Main(3),mk_n_list_Main 3,
        Length_Main, length_Main)) = 3)
```

## 5. Experiments (with demonstrations)

#### Example 1

```
class List { Int contain() { return assert false; } }
class Nil extends List { Int contain() { return false; } }
class Cons extends List {
    Int hd; List tl;
    Cons(Int hd, List tl) { this.hd=hd; this.tl=tl;}
    Int contain(Int n) {
        return if key = hd then true else this.tl.contain(key);
    }
}
main() {
    return assert (new Cons(1, new Cons(2, new Nil()))).contain(2);
```

Result: Safe (+ refinement types as certificate) Achieved fully-automated context-sensitive verification

Example 2

FJ program  $P_{list}$  + Annotation for recursive data types

Result: Safe (+ refinement types as certificate)

Achieved path- and context-sensitive verification with small annotation burden

### 6. Future work

Improve RCaml (improve performance, etc.)
Improve the translation (deal with other features such as multithreaded programs, assignment, exceptions, etc.)