

## ソフトウェア論理 Logic in Computer Software

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### Strings as codes (1)

A standard C-program for the power function (べき乗を求める関数):

```
int power (int n, int x) {
    if (n == 1) {
        return x;
    } else if (even(n)) {
        return sqr(power(n/2,x));
    } else {
        return x*power(n-1,x);
    }
}
```

How to represent programs (codes) as data ?

- Strings
- Data types for trees
- Language support for code generation (Built-in data types)

(To distinguish two kinds of programs from each other, we write “programs” for **generating** programs, and “codes” for **generated** programs.)

This choice greatly affects the quality of programs and codes. (ease of writing/understanding, reusability efficiency, reliability, etc.)

### Strings as codes (2)

Suppose  $n$  is known now, and  $x$  is not known now. A generator for the power function in C-like notation:

```
string gen_power1 (int n, string xs) {
    if (n == 1) { return xs;
    } else if (even(n)) {
        return concat("sqr(", gen_power1(n/2,xs), ")");
    } else {
        return concat(xs, "*(", gen_power1(n-1,xs), ")");
    }
}
string gen_power (int n) {
    return
        concat("int power (int x) { return(",
              gen_power1(n, "x"), ");}");
}
```

assuming that concat does the right job.

## Strings as codes (3)

Inner product of vectors in C-like notation:

```
float ip (int n, float a[], float b[]) {
    int i;
    float sum = 0.0;
    for (i = 0; i < n; i++) {
        sum += a[i] * b[i];
    }
    return sum;
}
```

## Strings as codes (4)

Suppose  $n$  is known,  $a$  and  $b$  are not known. Generator for inner product:

```
string gen_ip1 (int n, int idx, string as, string bs) {
    if (idx == n) return "0.0";
    else return
        concat(as, "[", int_to_string(idx), "] * ",
              bs, "[", int_to_string(idx), "] + ",
              gen_ip1(n, idx + 1, as, bs));
}
string gen_ip (int n, string as, string bs) {
    return
        concat("float ip (int ", as, "[], int", bs, "[]) {"
              "return ", gen_ip1(n, 0, as, bs), ";", "}");
}
```

## Strings as codes (5)

Sometimes, we want to generate more specialized code: Suppose  $n$  and  $a$  are known, and  $b$  is not known.

```
string gen_ip1 (int n, int idx, float a[], string bs) {
    if (idx == n) return "0.0";
    else return
        concat(float_to_string(a[idx]), " * ",
              bs, "[", int_to_string(idx), "] + ",
              gen_ip1(n, idx + 1, a, bs) );
}
string gen_ip (int n, float a[], string bs) {
    return concat("float ip (int ", bs, "[]) {"
                  "return ", gen_ip1(n, 0, a, bs), ";", "}");
}
```

## Strings as codes (summary)

Evaluation:

- (+) It can be done in almost all programming languages.
- (+) So, we don't have to learn more techniques.
- (-) It needs a certain (boring) rewriting of the non-generating version
- (-) It is error prone, especially when we embed codes into code ("splicing")
- (-) It is not composable; we cannot combine one generator with internal variables "x" and "y", and another generator with internal variables "x" and "z".
- (-) Sometimes (or, often) the generated codes cannot be compiled due to type errors or unbound variables.

## Data types for trees as codes (1)

Lisp/Scheme has trees as primitive data (“Symbolic expression” or S-expression for short):

```
(+ 1 2) returns 3
'+ 1 2) returns (+ 1 2)
(list (+ 1 2) (* 2 3)) returns 9
(list '+ 1 2) '* 2 3)) returns ((+ 1 2) (* 2 3))
```

Suitable for symbolic computation (mathematical formulas, logical formulas, programs, XML data, sentences in natural languages etc.)

## Data types for trees as codes (2)

Power function in Scheme:

```
(define (power n x)
  (if (= n 1) x
      (if (even n)
          (sqr (power (/ n 2) x))
          (* x (power (- n 1) x))))))
```

## Data types for trees as codes (3)

Generator for Power function in Scheme:

```
(define (gen_power1 n xs)
  (if (= n 1) xs
      (if (even n)
          (list 'sqr (gen_power1 (/ n 2) xs))
          (list '* xs (gen_power1 (- n 1) xs)))))
```

```
(define (gen_power n)
  (list 'define '(power x)
        (gen_power1 n 'x)))
```

Slightly better than the “strings as codes” approach.  
Still splicing is problematic.

## Data types for trees as codes (4)

(from the previous slide)

```
(define (gen_power n)
  (list 'define '(power x)
        (gen_power1 n 'x)))
```

Generator for Power function in Scheme using **quasi-quotation**:

```
(define (gen_power n)
  '(define (power x)
    ,(gen_power1 n 'x)))
```

Can represent splicing neatly.

Quasi-quotation is like quotation, but allows splicing.

## Data types for trees as codes (5)

Evaluation:

- (+) Better syntax. Ease of writing and understanding. Much less error-prone.
- (+) No overhead; runs in exactly the same speed as the one without quasi-quotation (it is just an input-macro).
- (-) Programming language (or its preprocessor) must support it.
- (-) Still not composable; we cannot combine one generator with internal variables "x" and "y", and another generator with internal variables "x" and "z".
- (-) Sometimes (or, often) the generated codes cannot be compiled due to unbound variables.

## Language support (built-in data types) (2)

Generator for Power:

```
let rec gen_power1 n xs =
  if n=1 then xs
  else if (even n) then
    '(sqr ,(gen_power1 (n / 2) xs))
  else '(,xs * ,(gen_power1 (n - 1) xs))
```

```
let gen_power n =
  '(fun x -> ,(gen_power1 n 'x))
```

## Language support (built-in data types) (1)

Power in OCaml (a dialect of ML):

```
let rec power n x =
  if n=1 then x
  else if (even n) then
    sqr (power (n / 2) x)
  else x * (power (n-1) )
```

## Language support (built-in data types) (2')

Generator for Power:

```
let rec gen_power1 n xs =
  if n = 1 then xs
  else if (even n) then
    <sqr ~(gen_power1 (n / 2) xs)>
  else <~xs * ~(gen_power1 (n - 1) xs)>
```

```
let gen_power n =
  <fun x -> ~(gen_power1 n <x>>
```

Intuitively: <a b c> is '(a b c) and <a ~b c> is '(a ,b c)

Then, we have:

```
gen_power 3 <x>
-> < ~<x> * ~(gen_power 2 <x>) >
-> < x * ~(<sqr ~(gen_power 1 <x>>)> >
-> < x * ~(<sqr ~(<x>>)> >
-> < x * ~(<sqr x>>
```

## Language support (built-in data types) (3)

But why is it better than Lisp/Scheme ?

Support for types.

- Types give a certain reliability of generator.
- Types give a certain reliability of generated codes,
- AND it ensures “no free variables” in generated codes.

Errors:

`x + 1`, `<x + 1>`, `<3.0 + 1>` `<~x + 1>`

Ok: `<fun x -> x + 1>`, `fun x -> <~x + 1>`,

`fun x -> <fun y-> ~x + y + 1>`,

## Summary

- “Codes as strings” are available in most languages, but no support.
- Staged computation: Language support for code generation.

## Language support (built-in data types) (4)

Type for codes

- If  $e$  is type  $\text{int}$ , then `< e >` is of type  $\text{int code}$ .
- In general, if  $e$  has type  $T$ , then `< e >` is of type  $T \text{ code}$ .
- If  $e$  has type  $T \text{ code}$ , then `~e` is of type  $T$ .

Types for `gen_power1`:

```
let rec gen_power1 n xs =
  if n = 1 then xs
  else if (even n) then
    .<sqr .~(gen_power1 (n / 2) xs)>.
  else .<~xs * .~(gen_power1 (n - 1) xs)>.
```

$n$  is of type  $\text{int}$ ,  $xs$  is of type  $\text{int code}$ .

the return type of the generator is  $\text{int code}$ .

then the generator has type

`int -> (int code) -> (int code)`.

## Exercises.

Assign types to the following terms (some terms do not have types.)

- `<fun x -> x+10> + 20` or `(+ '(lambda (x) (+ x 10)) 20)`
- `<fun x -> x+10> 20` or `('(lambda (x) (+ x 10)) 20)`
- `<(fun x -> x+10) 20>` or `'((lambda (x) (+ x 10)) 20)`
- `<fun x -> ~x + 10>` or `'(lambda (x) (+ ,x 10))`
- `fun y -> <fun x -> ~y +10>` or `(lambda (y) '(lambda (x) (+ ,y 10)))`
- `fun f -> <fun x -> ~(f <x+10>>>` or `(lambda (f) '(lambda (x) ,(f '(+ x 10))))`

## Exercises with Answers.

- `<fun x -> x+10> + 20` . (Not typable, since we cannot add 20 to a code, since a code itself is not an integer.)
- `<fun x -> x+10> 20` (Not typable, since we cannot apply a code to a value, since a code is not itself a function.)
- `<(fun x -> x+10) 20>` (Has a type `(int code)` under an empty typing context.)
- `<fun x -> ~x + 10>` (Has a type `(int code)` under a typing context `x :(int code)`.)
- `fun y -> <fun x -> ~y +10>` or (Has a type `int -> (int code)` under an empty typing context.)
- `fun f -> <fun x -> ~(f <x+10>>>` (Has a type `(int code -> int code) -> (int -> int) code` under an empty typing context.)